

# The largest real root of the independence polynomial of a unicyclic graph

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(Joint work with Iain Beaton)

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# Independence Polynomial

The **independence polynomial** of a graph  $G$ , denoted  $I(G, x)$ , is the generating polynomial for the number of independent sets of  $G$  of each order.

$$I(G, x) = \sum_{k=0}^{\alpha(G)} i_k^G x^k.$$

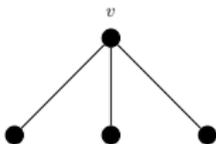
( $i_k^G$  denotes the number of independent sets of order  $k$ ,  $\alpha(G)$  denotes the independence number)

**Proposition (Gutman-Harary 1983):** If  $G$  is a graph and  $v \in V(G)$ , then

$$I(G, x) = I(G - v, x) + x \cdot I(G - N[v], x).$$

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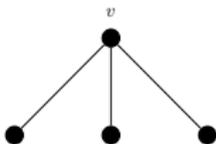


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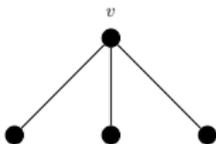


(b)  $S_4 - v$

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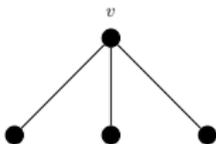


(b)  $S_4 - v$

$$I(S_4, x) = (1 + x)^3 + x \cdot 1$$

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(a)  $S_4$



(b)  $S_4 - v$

$$\begin{aligned} I(S_4, x) &= (1 + x)^3 + x \cdot 1 \\ &= 1 + 4x + 3x^2 + x^3 \end{aligned}$$

The roots of  $I(G, x)$  are called the **independence roots** of  $G$ .

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**Theorem (Brown-Hickman-Nowakowski 2004):** The set of all independence roots is dense in  $\mathbb{C}$ .

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Bounds on  $\beta$  are open for all other families of graphs.

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- 1)  $\preceq$  is transitive.
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**Theorem (Csikvári 2013)** If  $T$  is a tree of order  $n$ , then  $P_n \preceq T \preceq S_n$ .

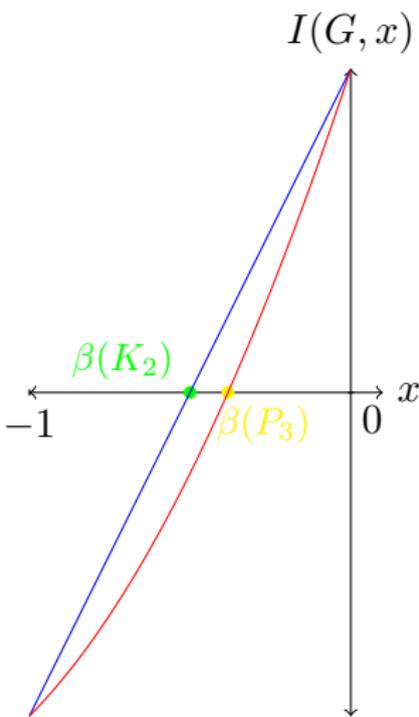


Figure:  $K_2 \preceq P_3$  so  $\beta(K_2) \leq \beta(P_3)$ .

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**Theorem (Oboudi 2018):** Let  $G$  and  $H$  be graphs with  $u \in V(G)$ ,  $e \in E(G)$ ,  $v \in V(H)$ , and  $f \in E(H)$ . Then the following hold:

- (i) If  $H - v \preceq G - u$  and  $G - N[u] \preceq H - N[v]$ , then  $H \preceq G$ .
- (ii) If  $H - f \preceq G - e$  and  $G - N[e] \preceq H - N[f]$ , then  $H \preceq G$ .

**Theorem (Beaton-C. 2020+):** If  $G$  is a connected unicyclic graph of order  $n$ , then  $C_n \preceq G \preceq U_n$ .

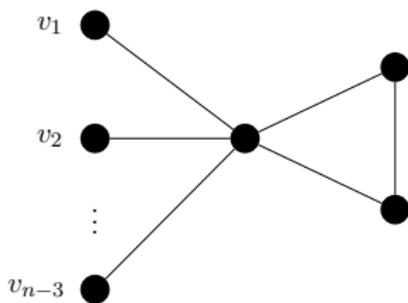


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**Corollary:** If  $G$  is a connected unicyclic graph of order  $n$ , then

$$\beta(C_n) \leq \beta(G) \leq \beta(U_n).$$

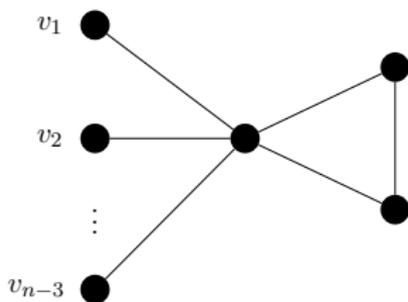


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**Theorem (Brown-Nowakowski. 2001):** Among all **well-covered** graphs of order  $n$ , the maximum modulus of an independence root is less than  $n$ .

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**Definition:** Let  $G$  be a graph. Form  $G^*$ , the **graph star** of  $G$  by attaching a leaf to each vertex of  $G$ .

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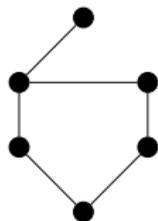


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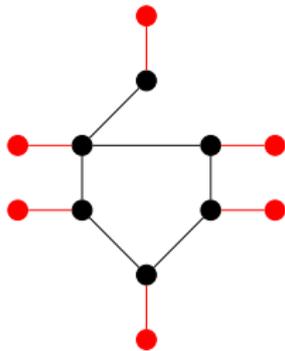


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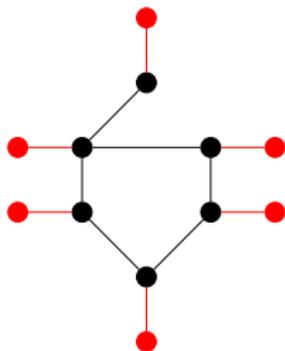


Figure:  $G^*$ .

**Theorem (Topp-Volkmann 1990):**  $G^*$  is well-covered for all graphs  $G$ .

**Theorem (Finbow-Hartnell-Nowakowski 1993):** If  $G \neq K_1$  and  $G \neq C_7$  and  $\text{girth}(G) \geq 6$ , then  $G$  is well-covered if and only if  $G = H^*$ .

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**Lemma (Beaton-C. 2020+):** Let  $G$  and  $H$  be graphs of order  $n$ . Then  $H \preceq G$  if and only if  $H^* \preceq G^*$ .

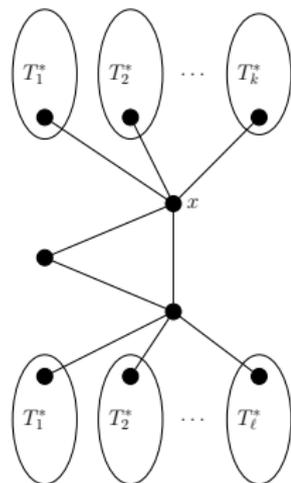
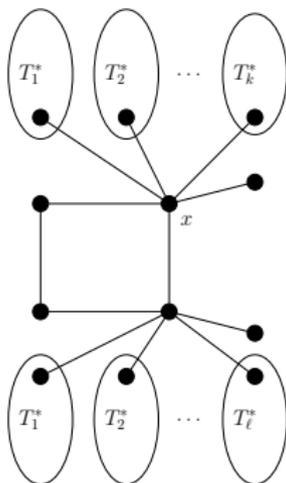
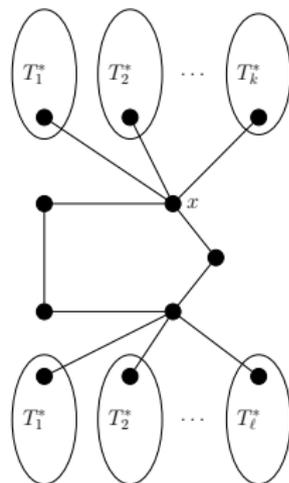
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**Corollary:** If  $T$  is a well-covered tree of order  $2n$ , then  $P_n^* \preceq T \preceq S_n^*$ .

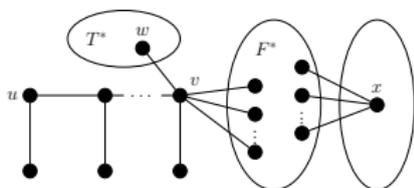
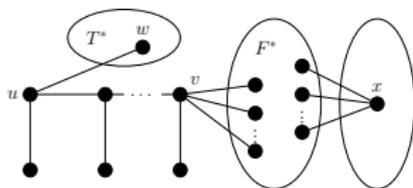
**Theorem (Topp-Volkmann 1990):** A graph  $G$  is a connected well-covered unicyclic graph if and only if

$$G \in \{C_3, C_4, C_5, C_7\} \cup \mathcal{S}_3 \cup \mathcal{S}_4 \cup \mathcal{S}_5 \cup \mathcal{KU}.$$

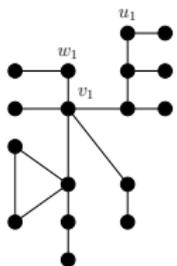
(a)  $\mathcal{S}_3$ (b)  $\mathcal{S}_4$ (c)  $\mathcal{S}_5$ 

**Note:** All graphs in  $\mathcal{S}_3$  and  $\mathcal{S}_5$  have odd order while all graphs in  $\mathcal{S}_4$  and  $\mathcal{KU}$  have even order by definition.

# The Dagger Operation

(a)  $G$ (b)  $G_{u,v,w}^{\dagger}$ 

**Lemma (Beaton-C. 2020+)** If  $G$  is a graph as above, then  $G_{u,v,w}^{\dagger} \preceq G$ .



(a)  $G_1$

Figure: Graph sequence formed by successive dagger operations.

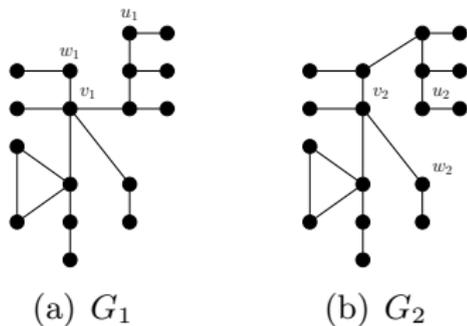


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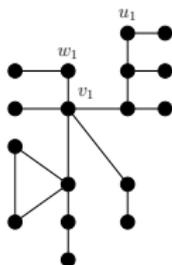
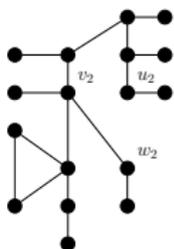
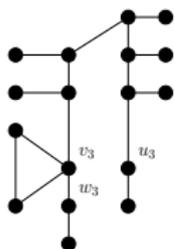
(a)  $G_1$ (b)  $G_2$ (c)  $G_3$ 

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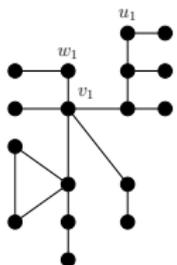
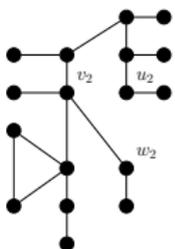
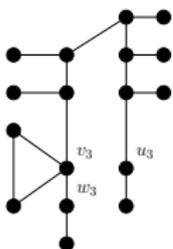
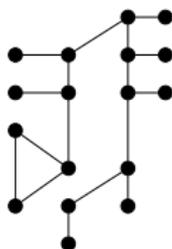
(a)  $G_1$ (b)  $G_2$ (c)  $G_3$ (d)  $G_4$ 

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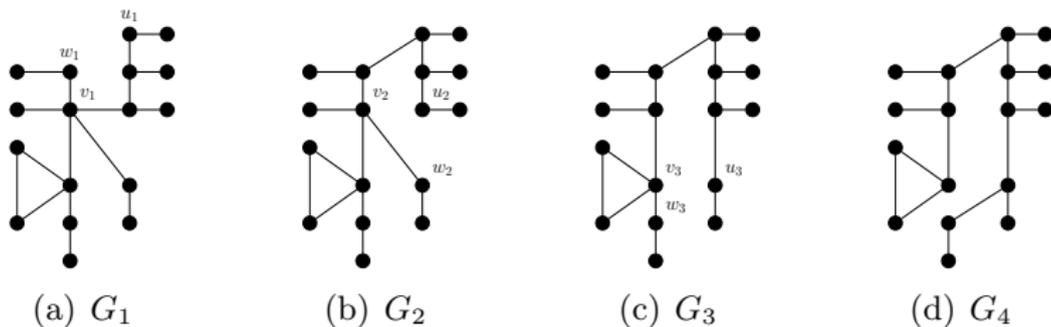


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Lemma (Beaton-C. 2020+): “This” works in general.

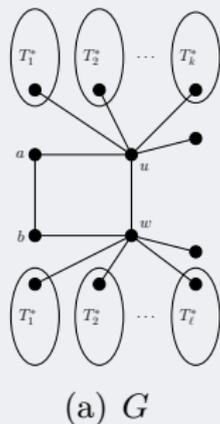
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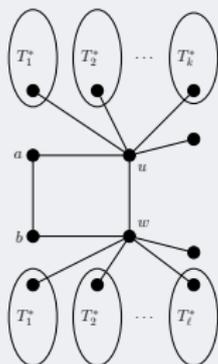
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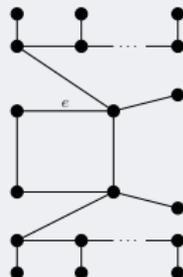
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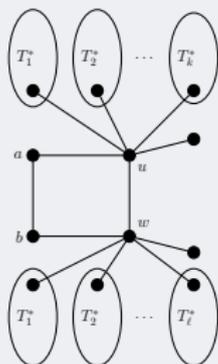
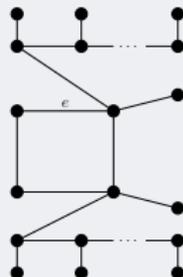
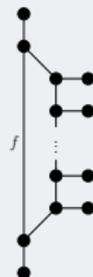


(b)  $F$

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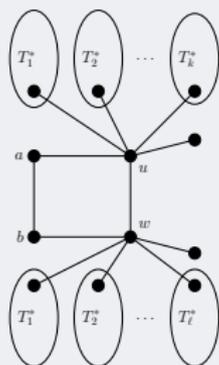
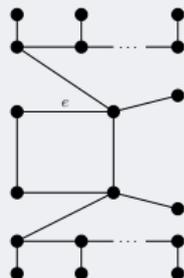
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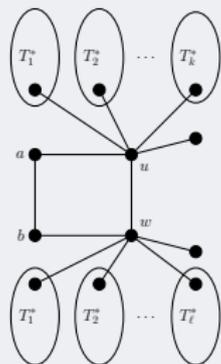
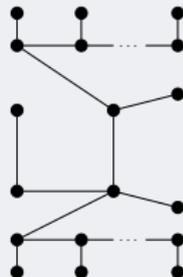
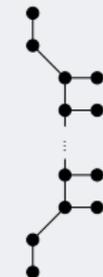
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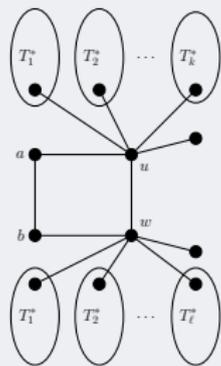
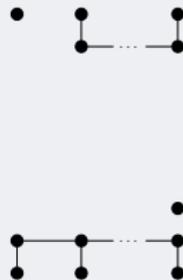
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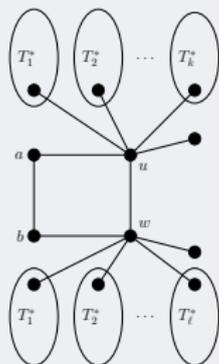
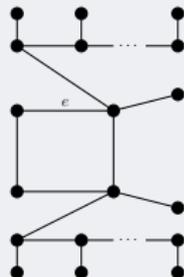
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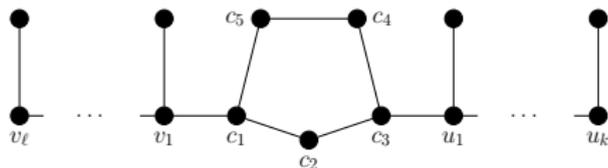
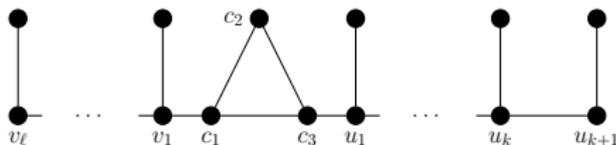
**Proof sketch lower bound.:** If  $G \in \mathcal{KU}$  done. Else,  $G \in \mathcal{S}_4$ .

(a)  $G$ (b)  $F$ (c)  $C_n^*$ 

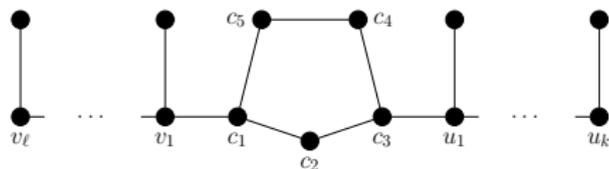
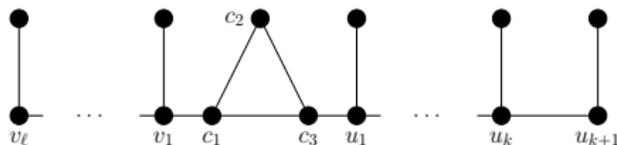
So  $G \succeq F \succeq C_n^*$ .



## Odd Order

(a)  $G(5, k, \ell)$ (b)  $G(3, k+1, \ell)$ Figure: Graphs in the family  $\mathcal{S}_P$ .

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**Lemma (Beaton-C. 2020+)** If  $G, H \in \mathcal{S}_P$  and have the same order, then  $I(G, x) = I(H, x)$ .

# Odd Order

**Theorem (Beaton-C. 2020+)** Let  $G$  be a connected well-covered unicyclic graph of odd order  $n$  and  $H_n \in \mathcal{S}_P$  of order  $n$ . Then

- i)  $C_n \preceq G \preceq M_n$  if  $n \leq 7$ , and
- ii)  $H_n \preceq G \preceq M_n$  if  $n \geq 9$ .

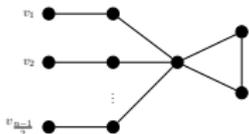


Figure: The graph  $M_n$ .

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**Corollary (Beaton-C. 2020+)** Let  $G$  be a connected well-covered unicyclic graph of odd order  $n$  and  $H_n \in \mathcal{S}_P$  of order  $n$ . Then

- i)  $\beta(C_n) \leq \beta(G) \leq \beta(M_n)$  if  $n \leq 7$ , and
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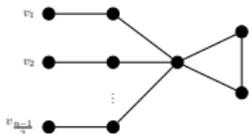


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**Conjecture:** If  $G$  is a triangle-free graph of order  $n$ , then  $P_n \preceq G \preceq K_{\lceil \frac{n}{2} \rceil, \lfloor \frac{n}{2} \rfloor}$ .

# THANK YOU!

